

Switching in the Rain: Predictive Wireless x-haul Network Reconfiguration

Igor Kadota*

Columbia University

Dror Jacoby*

Tel Aviv University

Hagit Messer

Tel Aviv University

Gil Zussman

Columbia University

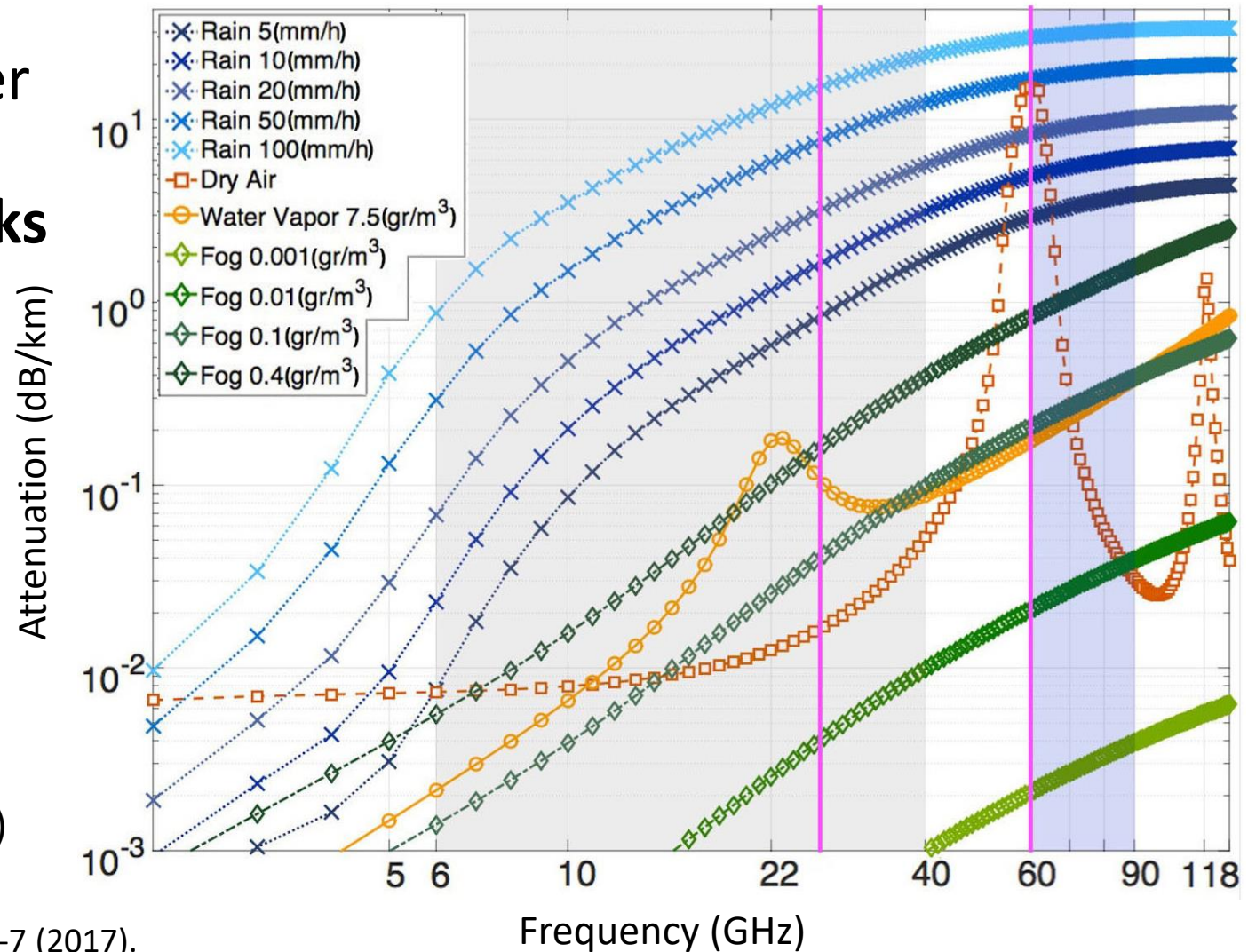
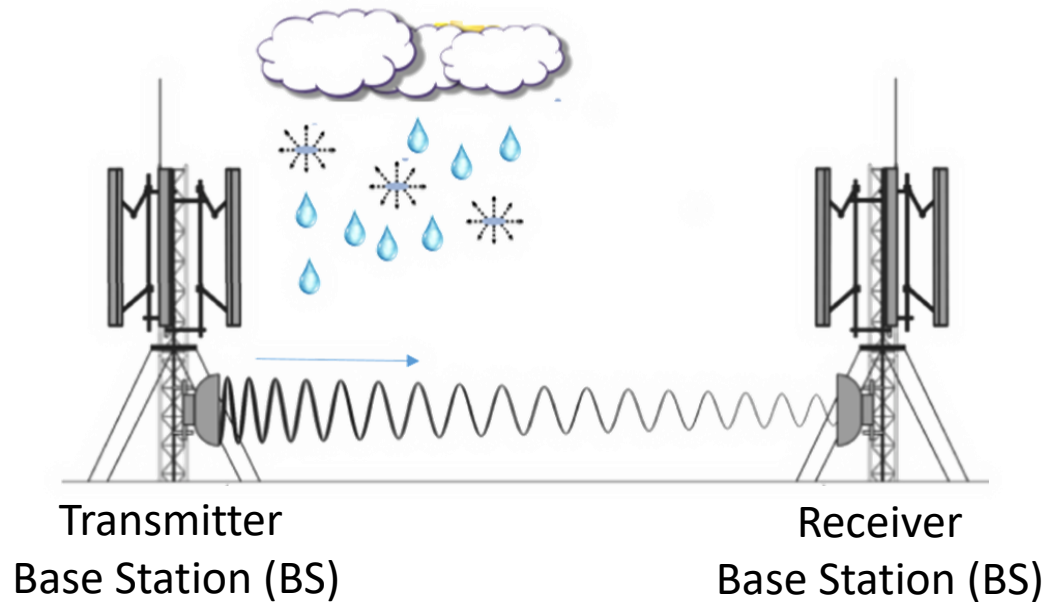
Jonatan Ostrometzky

Tel Aviv University

ACM SIGMETRICS, USA, June 19-22, 2023

Motivation

Different atmospheric and weather phenomena can cause **severe attenuation to high frequency links**



[21] ITU-R P.840. 2017. Attenuation due to clouds and fog. ITU 840-7 (2017).

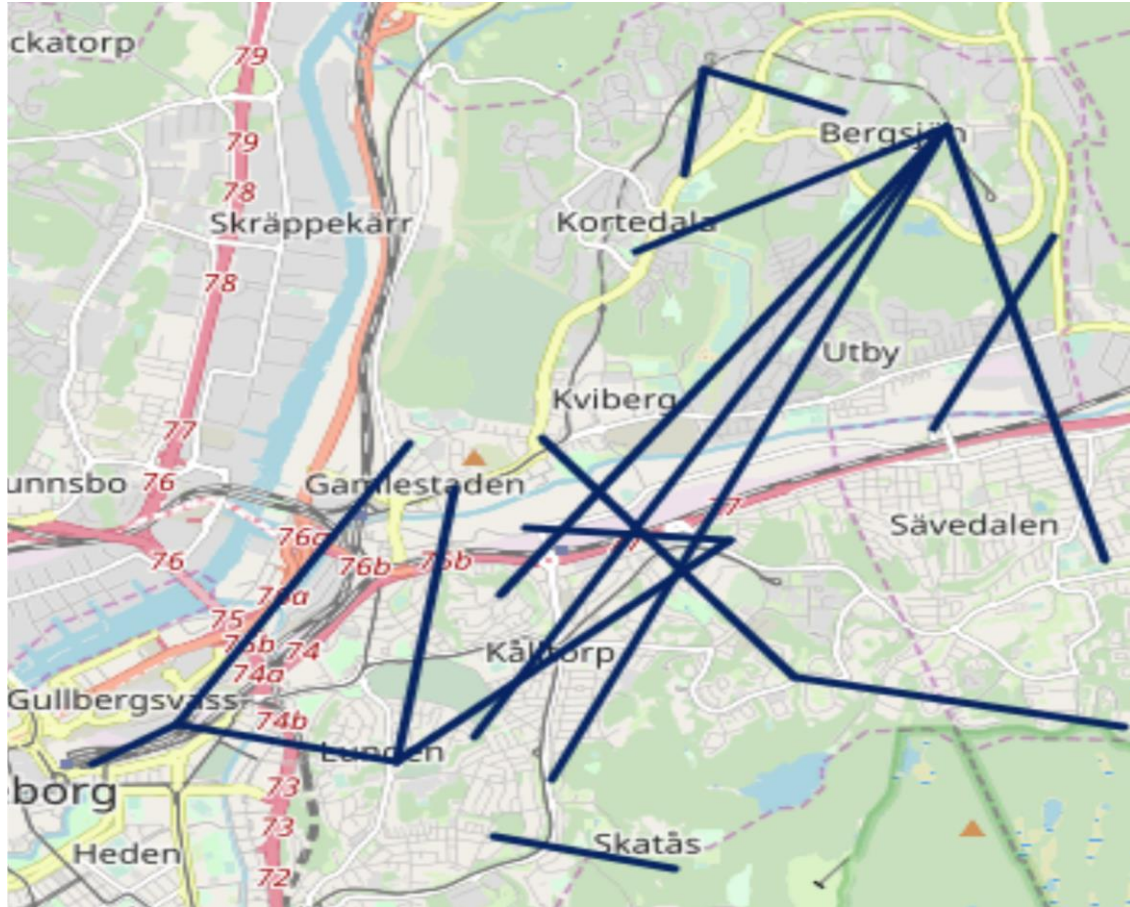
[22] ITU-R P.676. 2016. Attenuation by atmospheric gasses. ITU 676-11 (2016).

[23] ITU-R P.530. 2017. Propagation data and prediction methods required for the design of terrestrial line-of-sight systems. ITU 530-17 (2017).

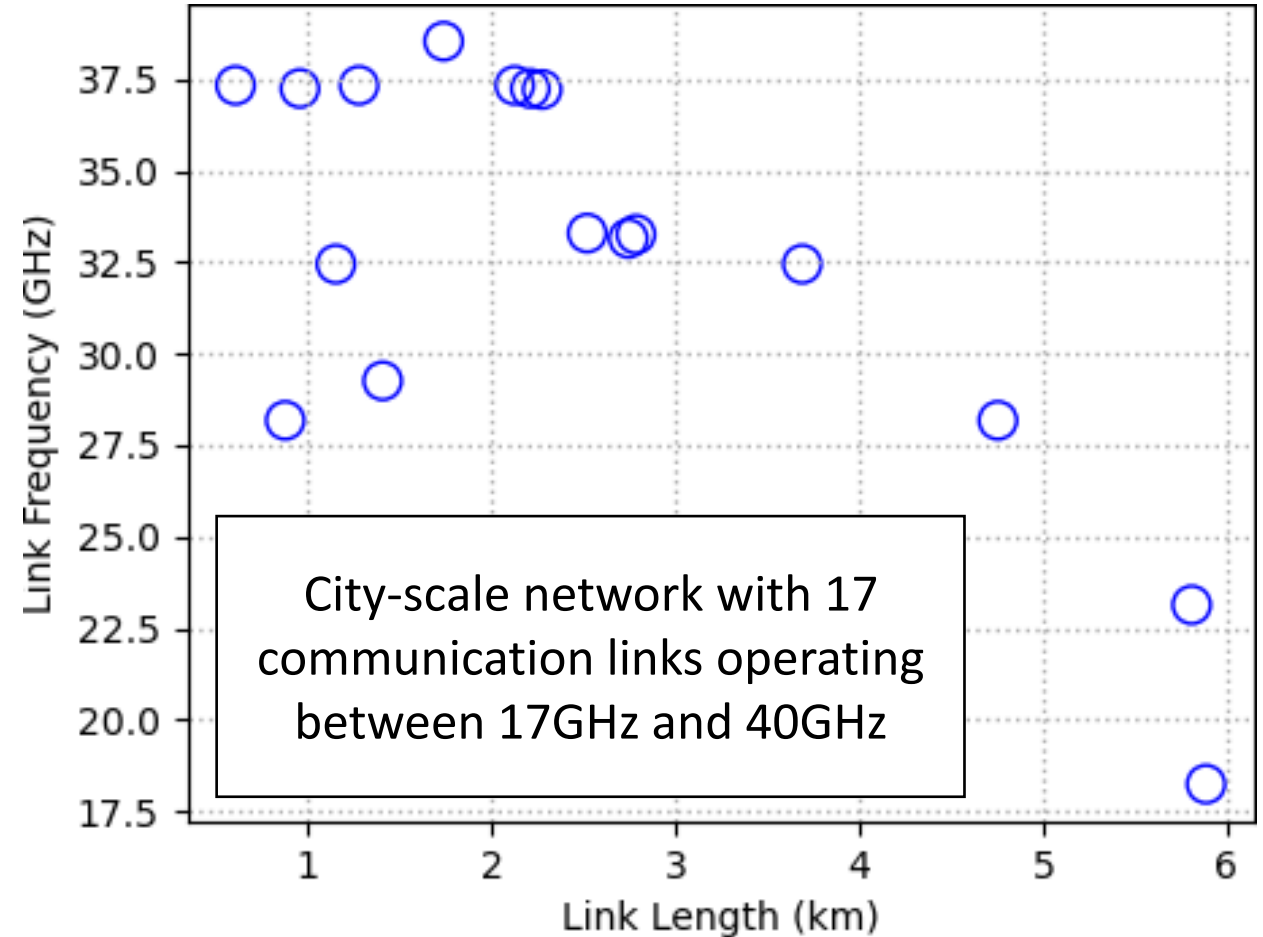
[24] ITU-R P.838. 2005. Specific attenuation model for rain for use in prediction methods. ITU 838-3, 1992-1999-2003-2005 (2005).

Motivation

Real-world city-scale wireless backhaul network in Gothenburg, Sweden

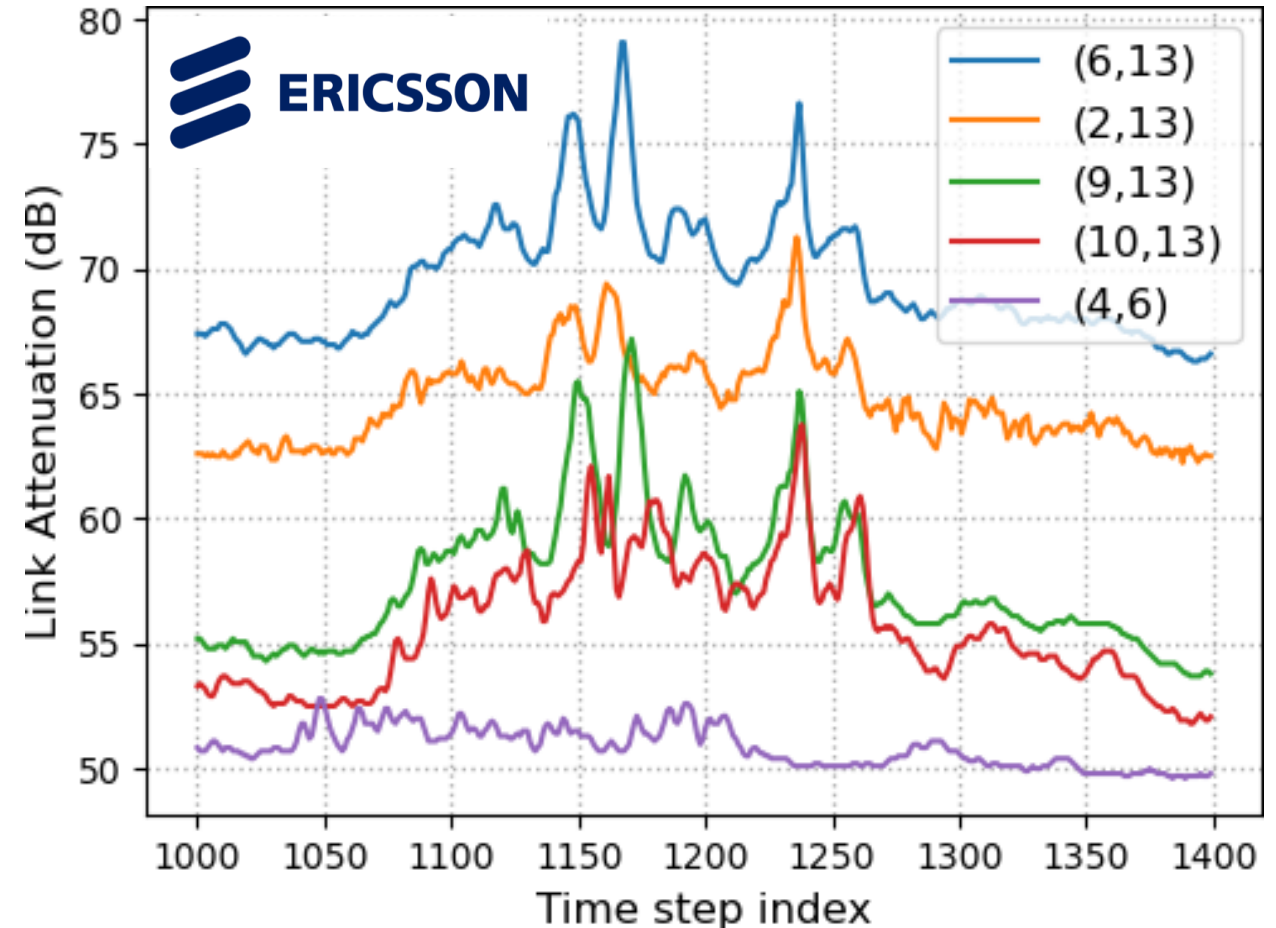
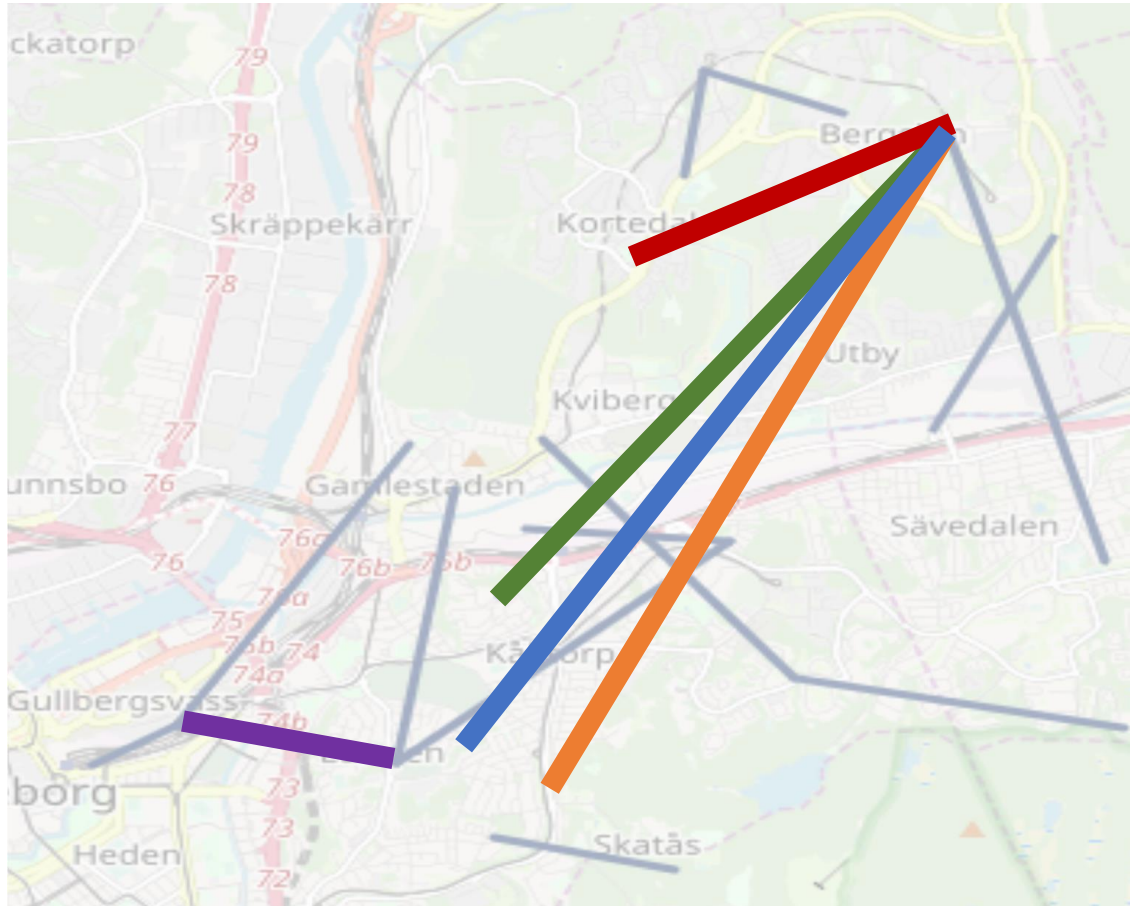


map area is of approximately 10x10 km²



Motivation

Ericsson collects link attenuation measurements [6] in time steps of 10sec



[6] L. Bao, C. Larsson, M. Mustafa, J. Selin, J. Andersson, J. Hansryd, M. Riedel, and H. Andersson. 2017. A brief description on measurement data from an operational microwave network in Gothenburg, Sweden. In Proc. CEST

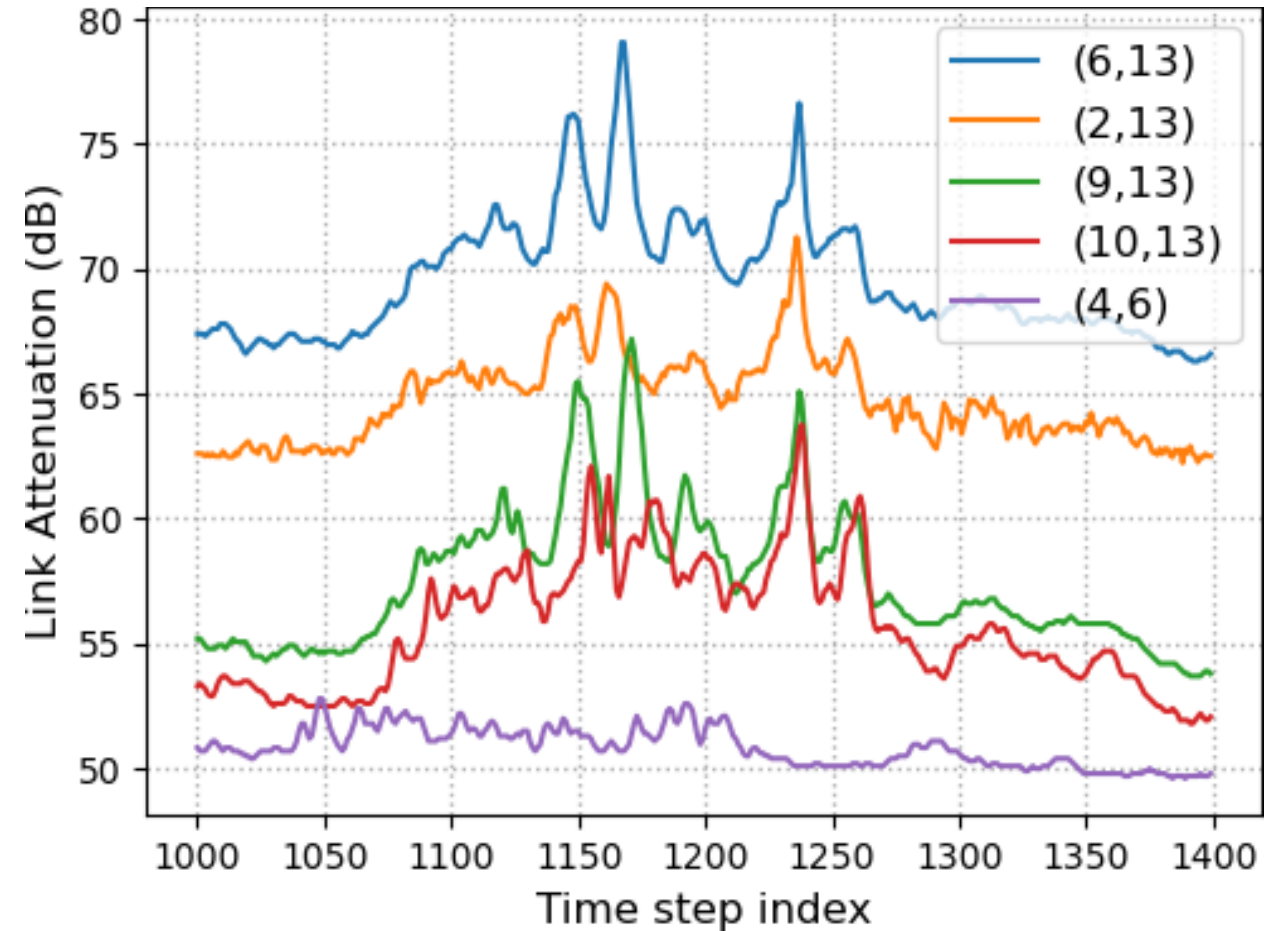
Motivation

Motivation: Need for a high-capacity wireless x-haul network that is **robust**

Challenge: **Link degradation varies** over time, geographical location, rain intensity, and **can be severe**.

Existing Solution:

- Global (**reactive**) NET layer mechanisms
 - NEC's SDN-based backhaul solution [40]



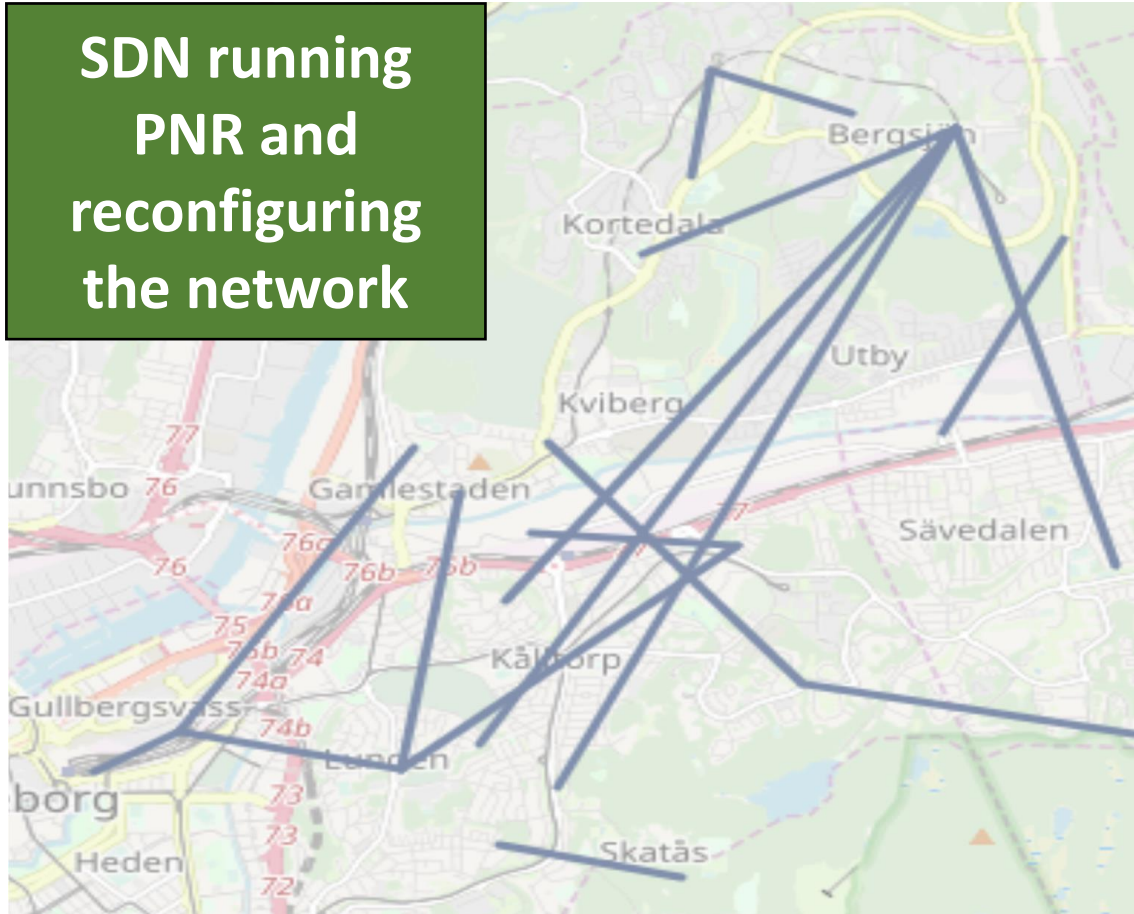
Related Work on Predictive Weather-Aware Reconfig.

Reference	Attenuation Prediction Mechanism			Routing Algorithm	Maximizes Throughput	Achieves Fair Allocation	Prevents Transient Congestion
	Weather Radar Measurements	Temporal Correlation	Spatial Correlation				
[25] 2009 A. Jabbar	X			Distributed (OSPF)			
[29] 2013 N. Javed		X		Distributed (OSPF)			
[45] 2016 J. Rak	X			Distributed (OSPF)			
[59] 2018 F. Yaghoubi		X		Centralized (SDN)	X		
This paper		X	X	Centralized (SDN)	X	X	X

J. Ostrometzky, G. Zussman, H. Messer, D. Jacoby, and I. Kadota. Predictive Weather-Aware Communication Network Management. US Patent Application No. 17/551,643. December 2021.

Predictive Network Reconfiguration (PNR)

SDN running
PNR and
reconfiguring
the network



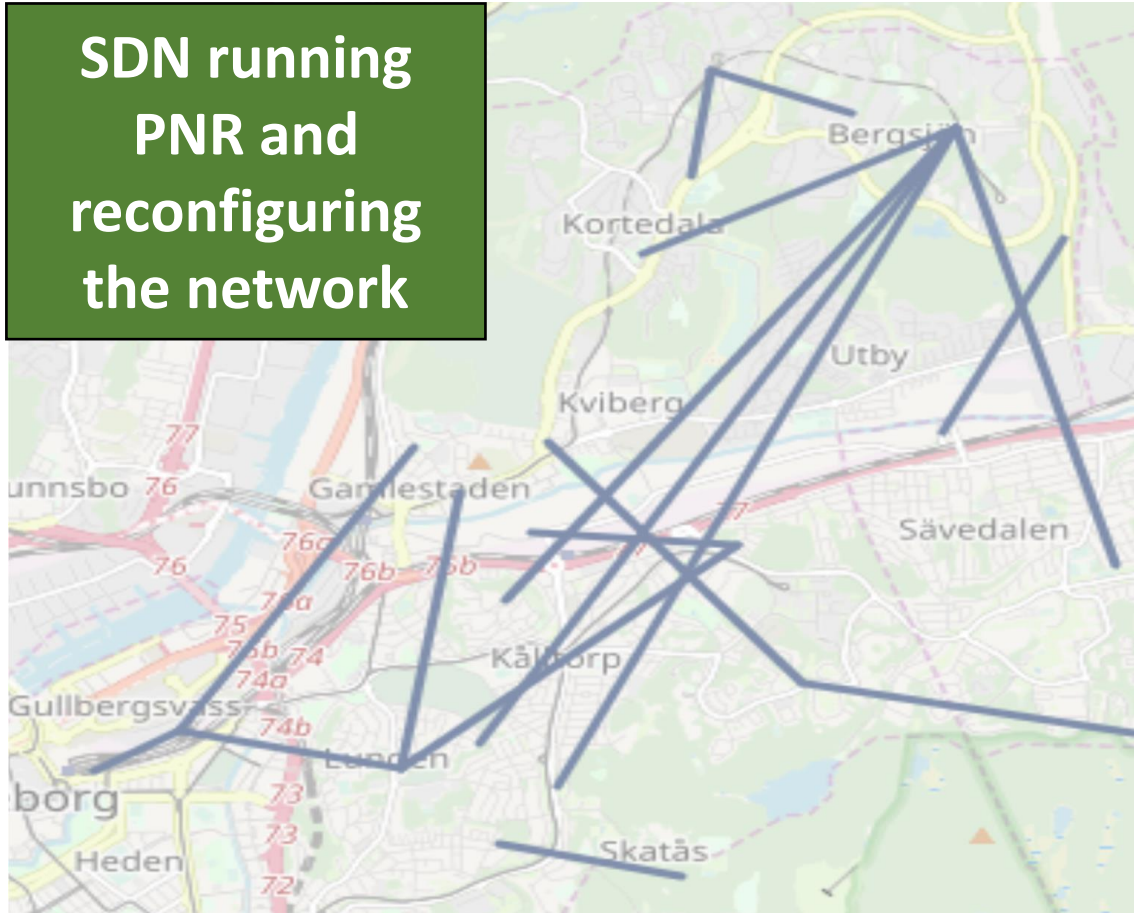
(Existing) Data Collection System [6]

Attenuation Prediction
(AP) Mechanism

Multi-Step Network Reconfiguration
(MSNR) Algorithm

Outline

**SDN running
PNR and
reconfiguring
the network**



(Existing) Data Collection System [6]

**Attenuation Prediction
(AP) Mechanism**

**Multi-Step Network Reconfiguration
(MSNR) Algorithm**

**Evaluation of the Predictive Network
Reconfiguration (PNR) framework**

Dataset

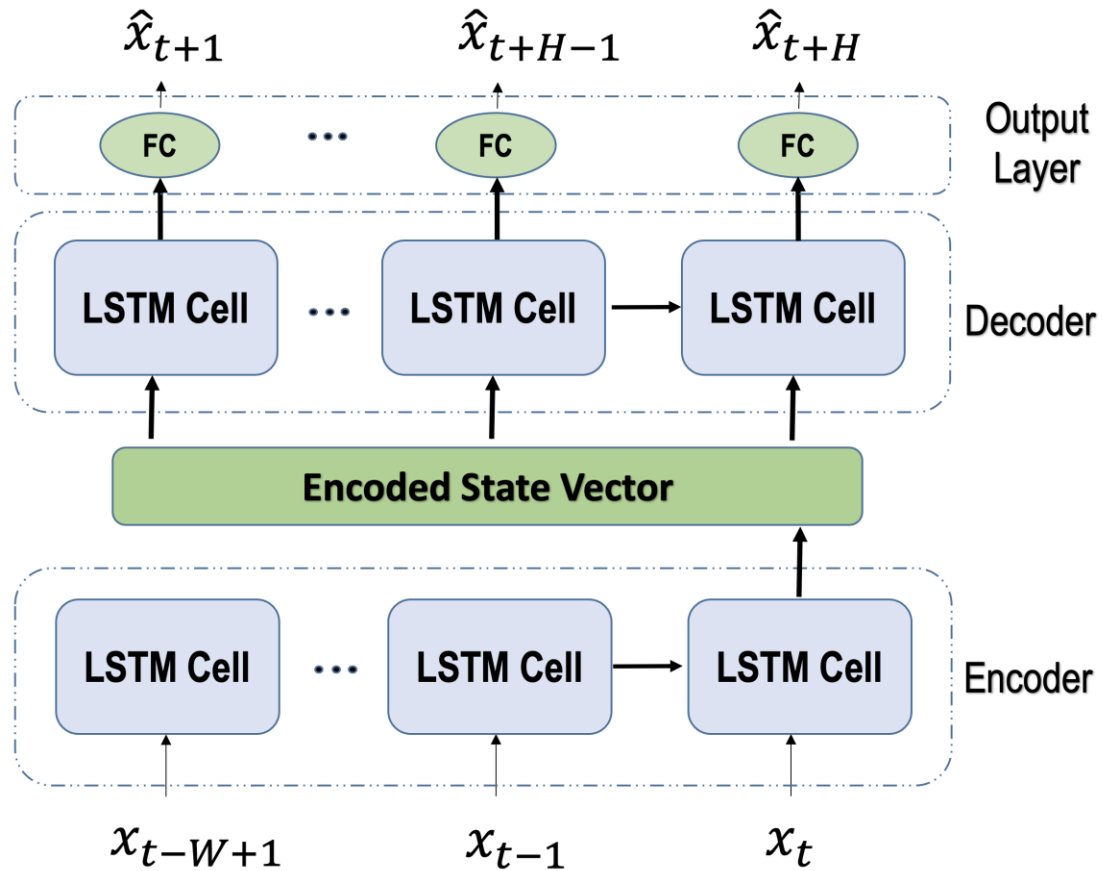


- 2,295,000 measurements of link attenuation levels $x_t^{(k,l)}$ for each of the 17 links, in intervals of $\Delta t = 10$ sec, containing both dry and rainy periods.

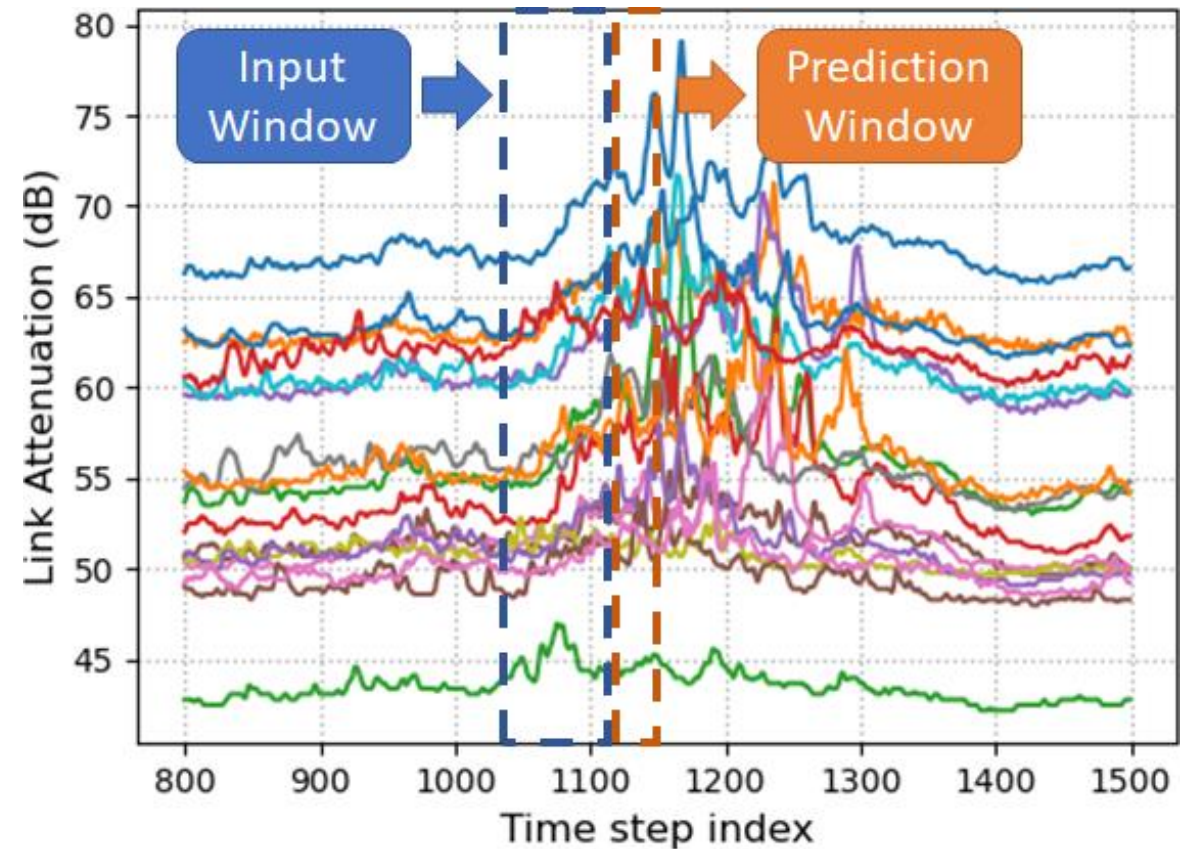


[3] J. Andersson, J. Olsson, R. van de Beek, and J. Hansryd. 2022. OpenMRG: Open data from Microwave links, Radar, and Gauges for rainfall quantification in Gothenburg, Sweden. *Earth System Science Data Discussions*.

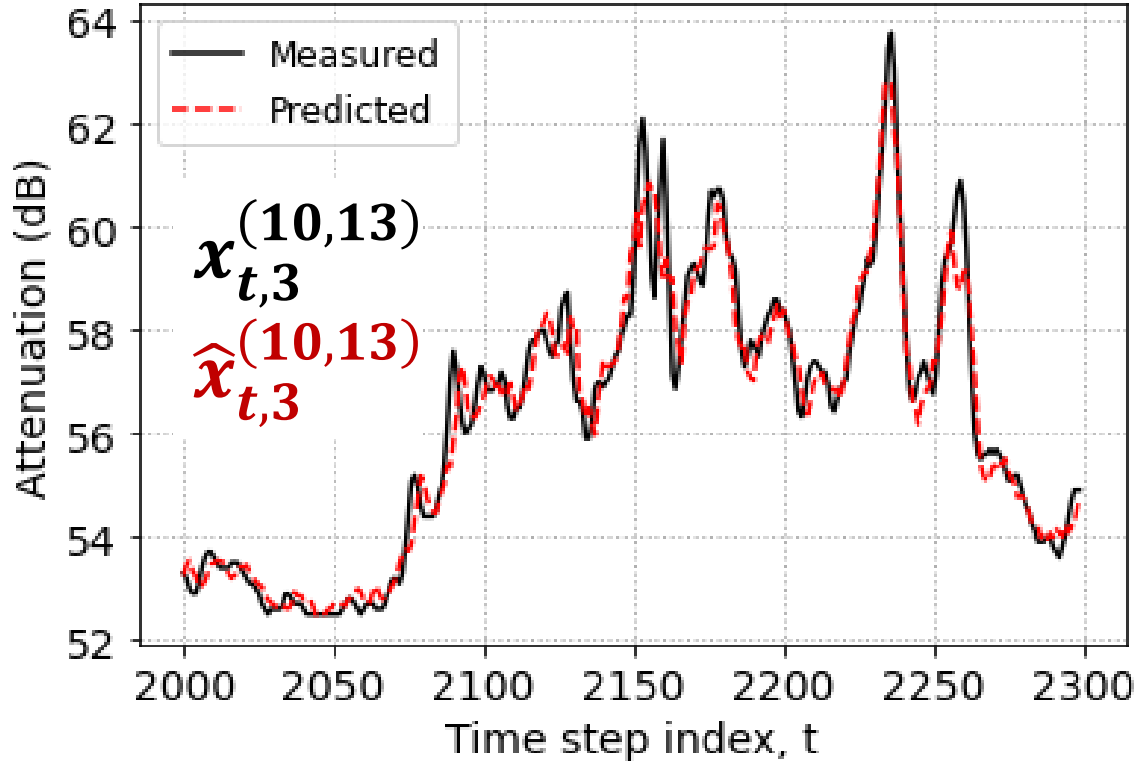
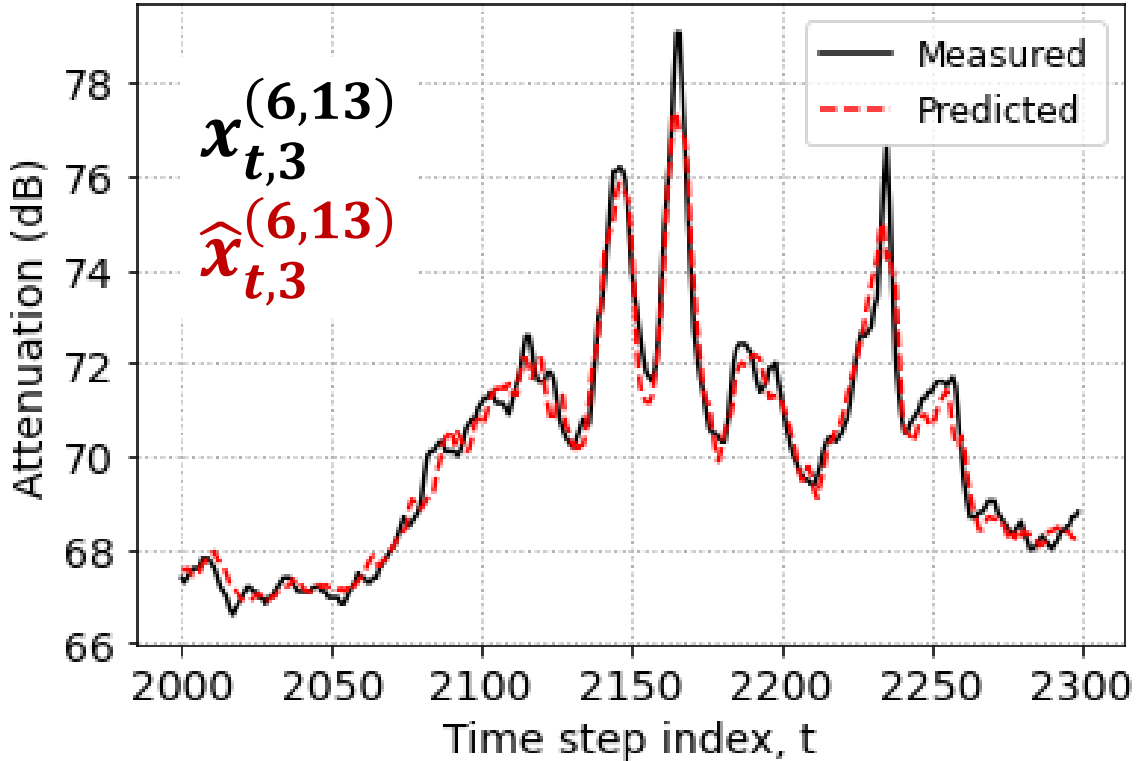
Encoder-Decoder LSTM Model



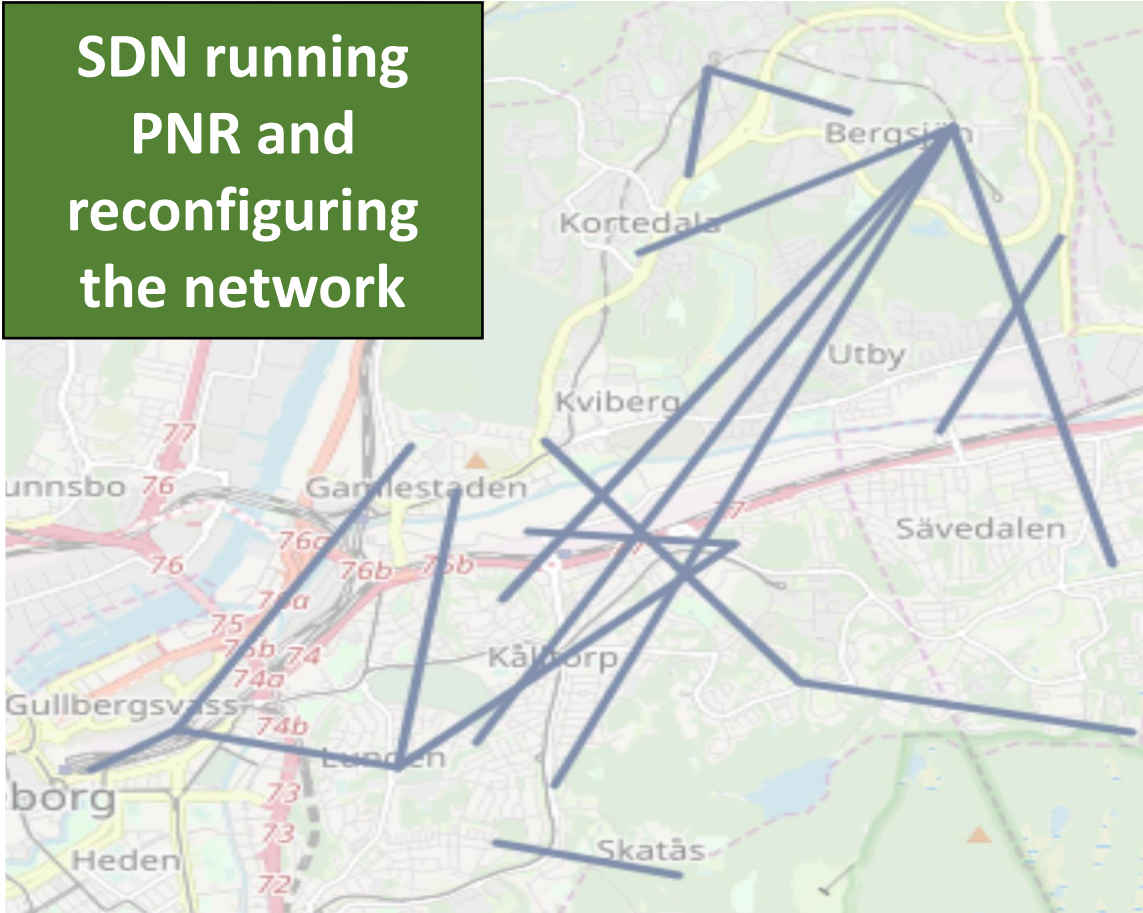
$$L(\theta) = \sum_{h=1}^H \sum_{\forall(k,l)} \left(x_{t+h}^{(k,l)} - \hat{x}_{t+h}^{(k,l)}(\theta) \right)^2$$



Evaluation of the AP Mechanism



Outline



(Existing) Data Collection System

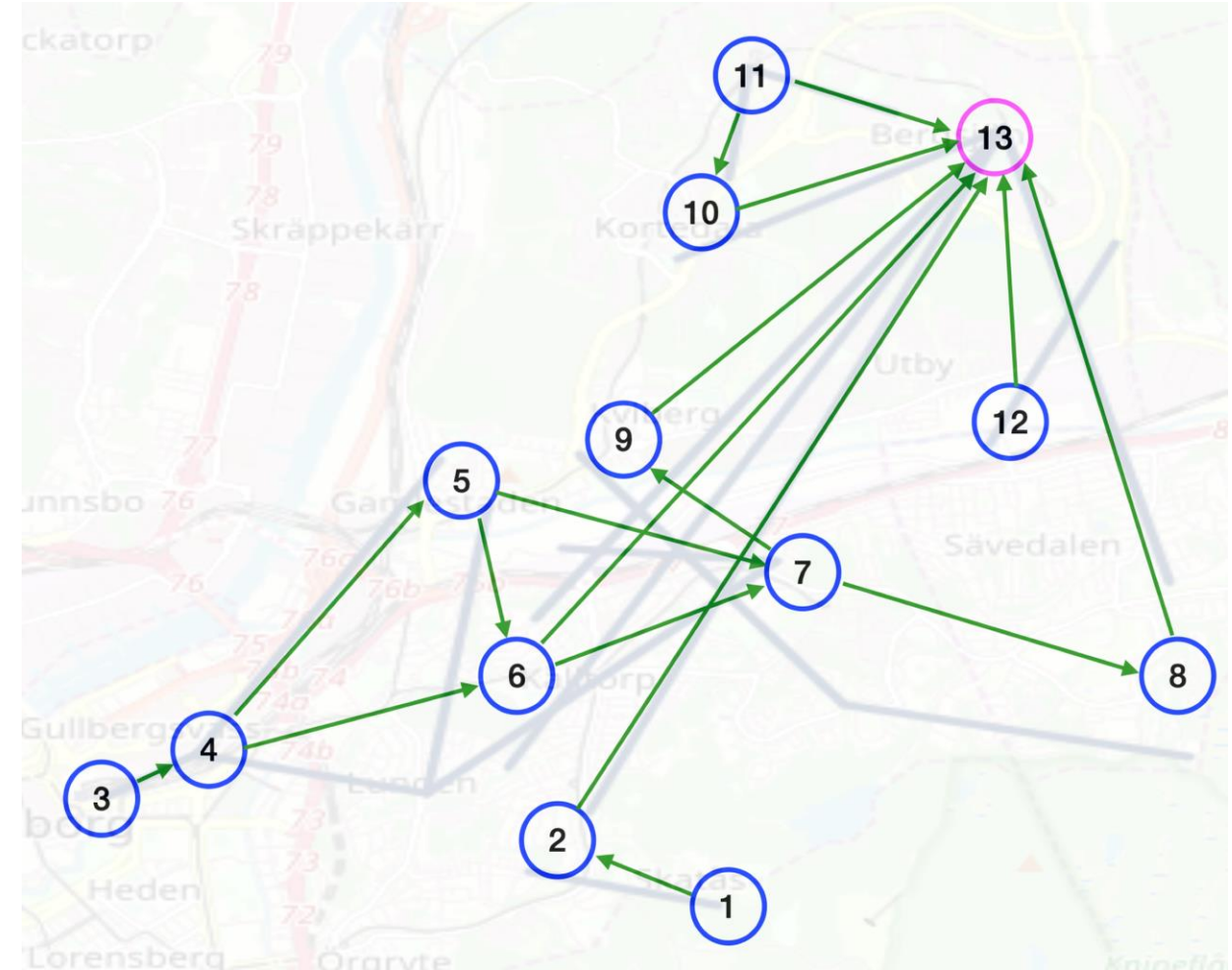
Attenuation Prediction (AP) Mechanism

Multi-Step Network Reconfiguration (MSNR) Algorithm

Evaluation of the Predictive Network Reconfiguration (PNR) framework

Network Model: Assumptions

- Neighboring link endpoints (up to 300m apart) form a Base Station (BS), also called a node
- Links are unidirectional
- Single destination BS (node $n = 13$)
- Time is divided in time-steps with index $t \in \{1, \dots, T\}$ and $\Delta t = 10\text{sec}$
- Admission Control and Routing decisions remain fixed between time-steps, e.g., t and $t + 1$.



Multi-Step Network Reconfiguration Problem

Goal: dynamically optimize **admission control** and **routing** decisions, i.e.:

$$\underset{\mathbf{z}_{n,t}, f_{n,t}^{(k,l)}}{\text{maximize}} \sum_{t=1}^T \sum_{n=1}^{N-1} \mathbf{z}_{n,t}$$

such that:

$$\sum_{n=1}^{N-1} \mathbf{z}_{n,t} d_n f_{n,t}^{(k,l)} \leq \min \left\{ c_t^{(k,l)}, \hat{c}_{t+1}^{(k,l)} \right\}, \forall (k,l) \in E, \forall t$$

$$\sum_{k=1}^N f_{n,t}^{(k,l)} + \mathbb{I}_{\{l=n\}} = \sum_{m=1}^N f_{n,t}^{(l,m)} + \mathbb{I}_{\{l=N\}}, \forall n \in V, \forall l \in V, \forall t$$

while guaranteeing that, **at every time-step t** , the selected $\left\{ \mathbf{z}_{n,t}, f_{n,t}^{(k,l)} \right\}$:

- achieves **max-min fairness** among the BSs sharing the network
- can be implemented **without inducing transient congestion**

Multi-Step Network Reconfiguration Problem

Max-min fairness: in order to maintain feasibility, $\uparrow \mathbf{z}_{n,t}$ necessarily results in $\downarrow \mathbf{z}_{m,t}$ for which $\mathbf{z}_{m,t} \leq \mathbf{z}_{n,t}$.

Transient congestion occurs when going from $\{\mathbf{z}_{n,t-1}, \mathbf{f}_{n,t-1}^{(k,l)}\}$ to $\{\mathbf{z}_{n,t}, \mathbf{f}_{n,t}^{(k,l)}\}$ results in the **violation of the capacity constraint of at least one link**.

while guaranteeing that, **at every time-step t** , the selected $\{\mathbf{z}_{n,t}, \mathbf{f}_{n,t}^{(k,l)}\}$:

- achieves **max-min fairness** among the BSs sharing the network
- can be implemented **without inducing transient congestion**

Known Solution to the **Static Single-Step** Problem

Goal: for a fixed t , optimize **admission control** and **routing** decisions, i.e.:

$$\underset{\mathbf{z}_{n,t}, \mathbf{f}_{n,t}^{(k,l)}}{\text{maximize}} \sum_{n=1}^{N-1} \mathbf{z}_{n,t}$$

such that:

$$\sum_{n=1}^{N-1} \mathbf{z}_{n,t} d_n \mathbf{f}_{n,t}^{(k,l)} \leq c_t^{(k,l)}, \forall (k,l) \in E$$

$$\sum_{k=1}^N \mathbf{f}_{n,t}^{(k,l)} - \sum_{m=1}^N \mathbf{f}_{n,t}^{(l,m)} = \mathbb{I}_{\{l=N\}} - \mathbb{I}_{\{l=n\}}, \forall n \in V, \forall l \in V$$

while guaranteeing that, the selected $\{\mathbf{z}_{n,t}, \mathbf{f}_{n,t}^{(k,l)}\}$:

- achieves **max-min fairness** among the BSs sharing the network

[2] M. Allalouf and Y. Shavitt. 2008. Centralized and distributed algorithms for routing and weighted max-min fair bandwidth allocation. *IEEE Trans. Netw. Service Manag.* 16, 5 (2008), 1015–1024.

[49] F. Shahrokhi and D. Matula. 1990. The maximum concurrent flow problem. *Journal of the ACM* 37, 2 (1990), 318–334.

Known Solution to the **Static Single-Step** Problem

Intuition of the solution [2,49] to the Maximum Concurrent Multi-Commodity Flow Problem:

- Iterative progressive filling ($\mathbf{z}_{n,t} \leftarrow \mathbf{z}$ for all unsaturated) & saturation test.

Potential Multi-Step Solution: at every time t , given the measured $c_t^{(k,l)}$, use the algorithm in [2] to obtain $\{\mathbf{z}_{n,t}, \mathbf{f}_{n,t}^{(k,l)}\}$.

Main Challenges:

- Reactive reconfiguration \rightarrow Reconfigure after impairment
- Transient congestion

[2] M. Allalouf and Y. Shavitt. 2008. Centralized and distributed algorithms for routing and weighted max-min fair bandwidth allocation. *IEEE Trans. Netw. Service Manag.* 16, 5 (2008), 1015–1024.

[49] F. Shahrokhi and D. Matula. 1990. The maximum concurrent flow problem. *Journal of the ACM* 37, 2 (1990), 318–334.

Proposed **Multi-Step** Algorithm: Key Elements

Cost of re-routing:

- to **guarantee zero transient congestion**, the SDN needs **scratch capacity**.
- at time t , when **not planning** to re-route at $t + 1$, the optimization should use the constraint : $\sum_{n=1}^{N-1} \mathbf{z}_{n,t} d_n \mathbf{f}_{n,t}^{(k,l)} \leq \min \left\{ c_t^{(k,l)}, \hat{c}_{t+1}^{(k,l)} \right\}, \forall (k, l) \in E$
- at time t , when **planning** to re-route at $t + 1$, the optimization should use the constraint: $\sum_{n=1}^{N-1} \mathbf{z}_{n,t} d_n \mathbf{f}_{n,t}^{(k,l)} \leq \min \left\{ c_t^{(k,l)}, \underbrace{(\mathbf{1} - \mathbf{s}) \hat{c}_{t+1}^{(k,l)}}_{\text{Cost of re-routing}} \right\}, \forall (k, l) \in E$

Proposed **Multi-Step** Algorithm: Key Elements

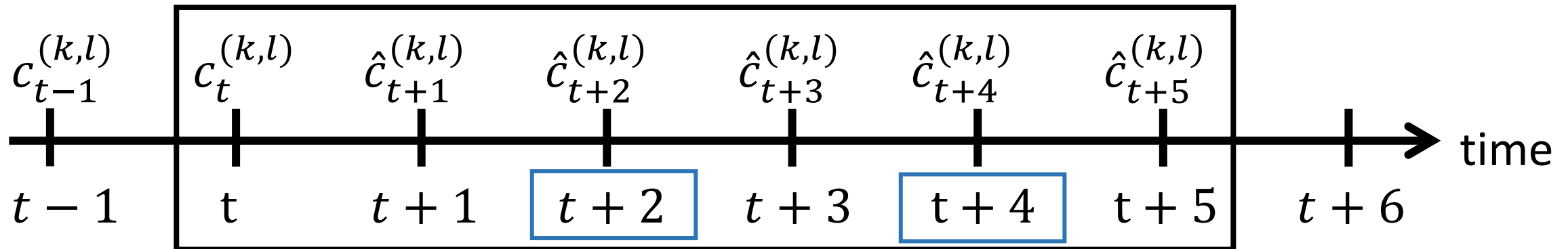
Joint Optimization over Multiple Time-Steps:

- Input: measured and predicted capacities $\{c_t^{(k,l)}, \hat{c}_{t+1}^{(k,l)}, \dots, \hat{c}_{t+H}^{(k,l)}\}$
- Output: sequence of $\{\mathbf{z}_{n,t}, f_{n,t}^{(k,l)}\}$ that maximizes $\sum_{h=0}^H \sum_{n=1}^{N-1} \mathbf{z}_{n,t+h}$

Proposed **Multi-Step** Algorithm: Key Elements

Joint Optimization over Multiple Time-Steps:

- Input: measured and predicted capacities $\{c_t^{(k,l)}, \hat{c}_{t+1}^{(k,l)}, \dots, \hat{c}_{t+H}^{(k,l)}\}$
- Output: sequence of $\{z_{n,t}, f_{n,t}^{(k,l)}\}$ that maximizes $\sum_{h=0}^H \sum_{n=1}^{N-1} z_{n,t+h}$
- MPC-based solution: consider every possible “**re-routing plan**” and select the plan that maximizes $\sum_{h=0}^H \sum_{n=1}^{N-1} z_{n,t+h}$. There are $\geq 2^H$ possible plans.
- One of these plans is illustrated below for $H = 5$:

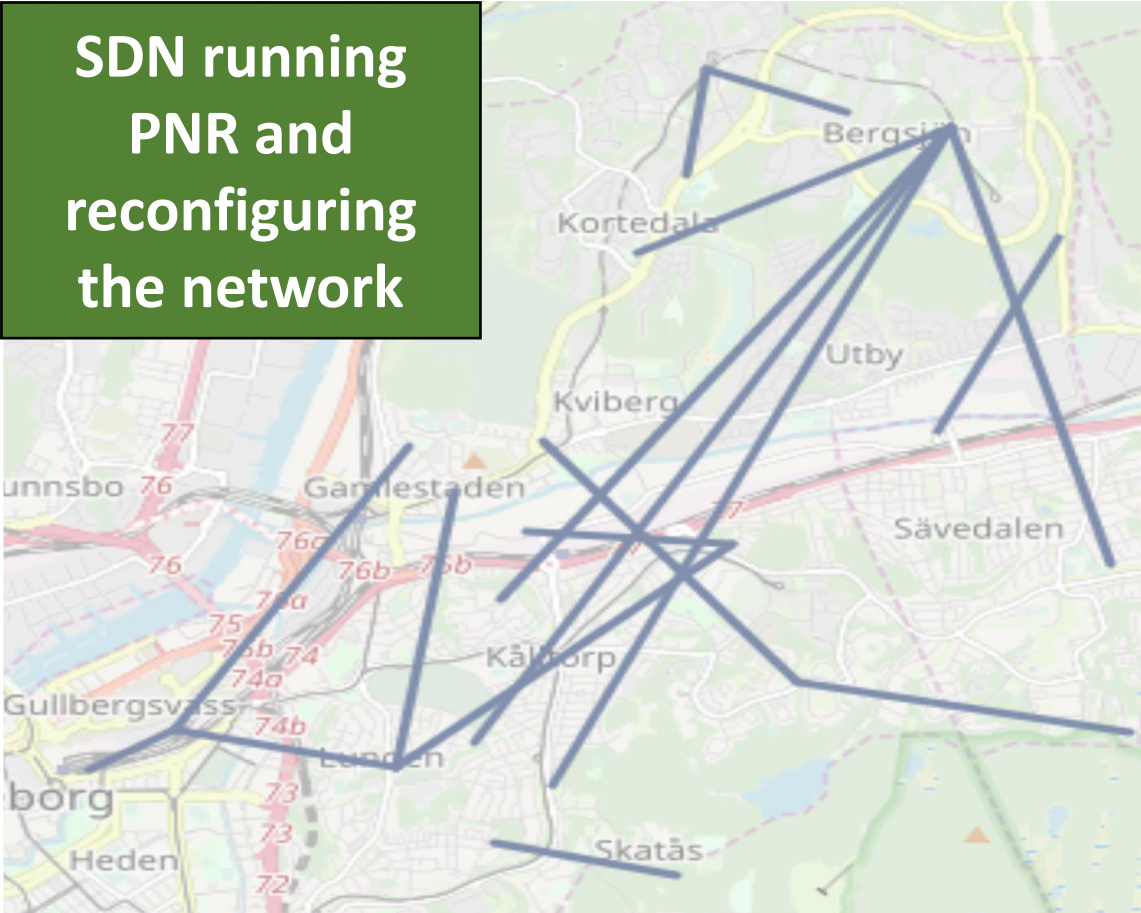


Proposed **Multi-Step** Algorithm: Key Elements

Joint Optimization over Multiple Time-Steps:

- Input: measured and predicted capacities $\{c_t^{(k,l)}, \hat{c}_{t+1}^{(k,l)}, \dots, \hat{c}_{t+H}^{(k,l)}\}$
- Output: sequence of $\{\mathbf{z}_{n,t}, \mathbf{f}_{n,t}^{(k,l)}\}$ that maximizes $\sum_{h=0}^H \sum_{n=1}^{N-1} \mathbf{z}_{n,t+h}$
- MPC-based solution: consider every possible “**re-routing plan**” and select the plan that maximizes $\sum_{h=0}^H \sum_{n=1}^{N-1} \mathbf{z}_{n,t+h}$. There are $\geq 2^H$ possible plans.
- Remark: using backward induction, we can reduce complexity from $O(2^H)$ to $O(H^4)$.
- Proposition: $\{\mathbf{z}_{n,t+h}, \mathbf{f}_{n,t+h}^{(k,l)}\}$ is max-min fair $\forall h$

Outline



(Existing) Data Collection System

Attenuation Prediction (AP) Mechanism

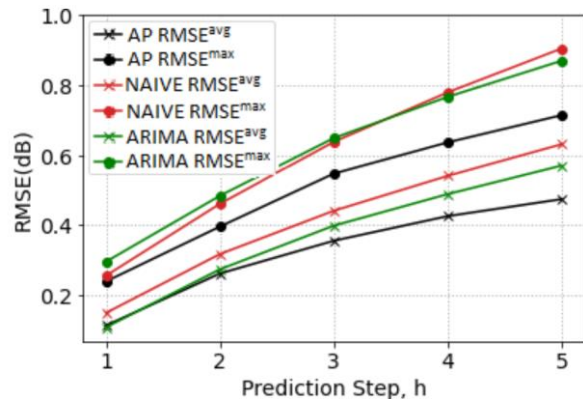
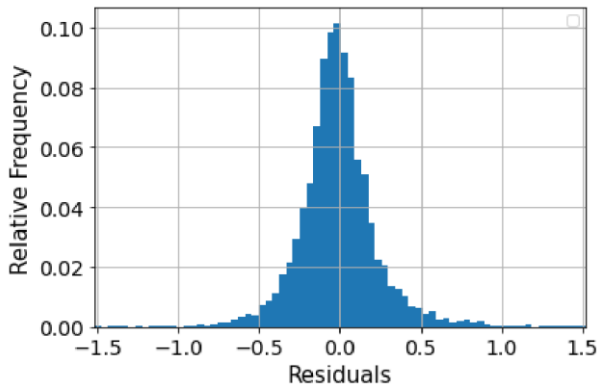
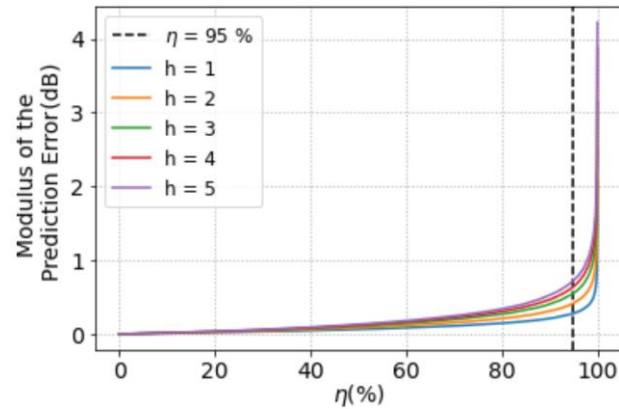
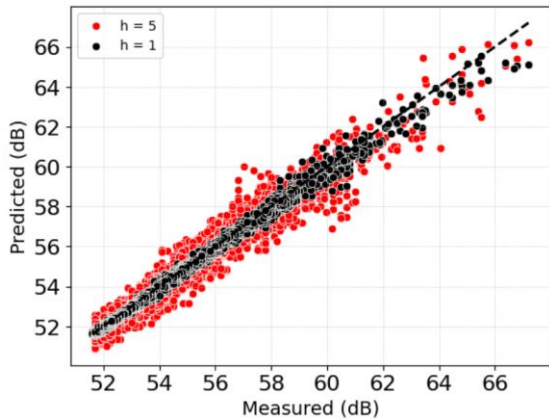
Multi-Step Network Reconfiguration (MSNR) Algorithm

Evaluation of the Predictive Network Reconfiguration (PNR) framework

Performance Evaluation

Attenuation Prediction

- Comparison with Naïve and ARIMA



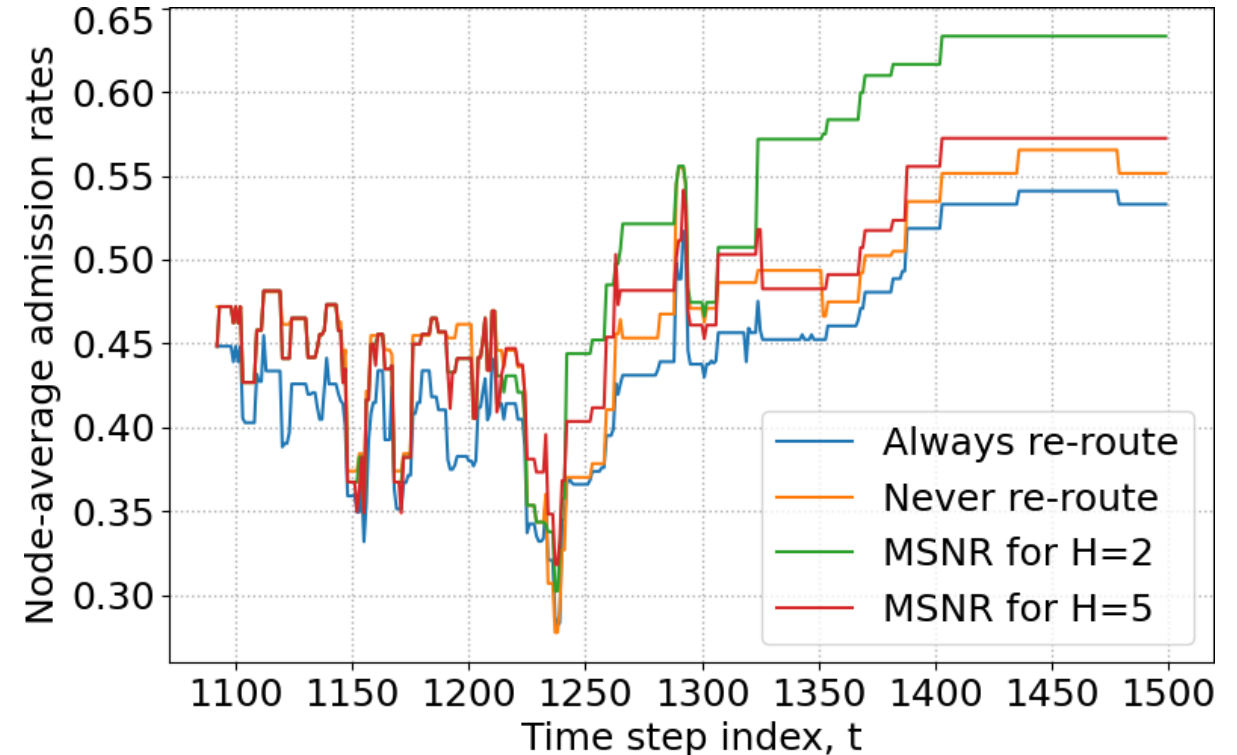
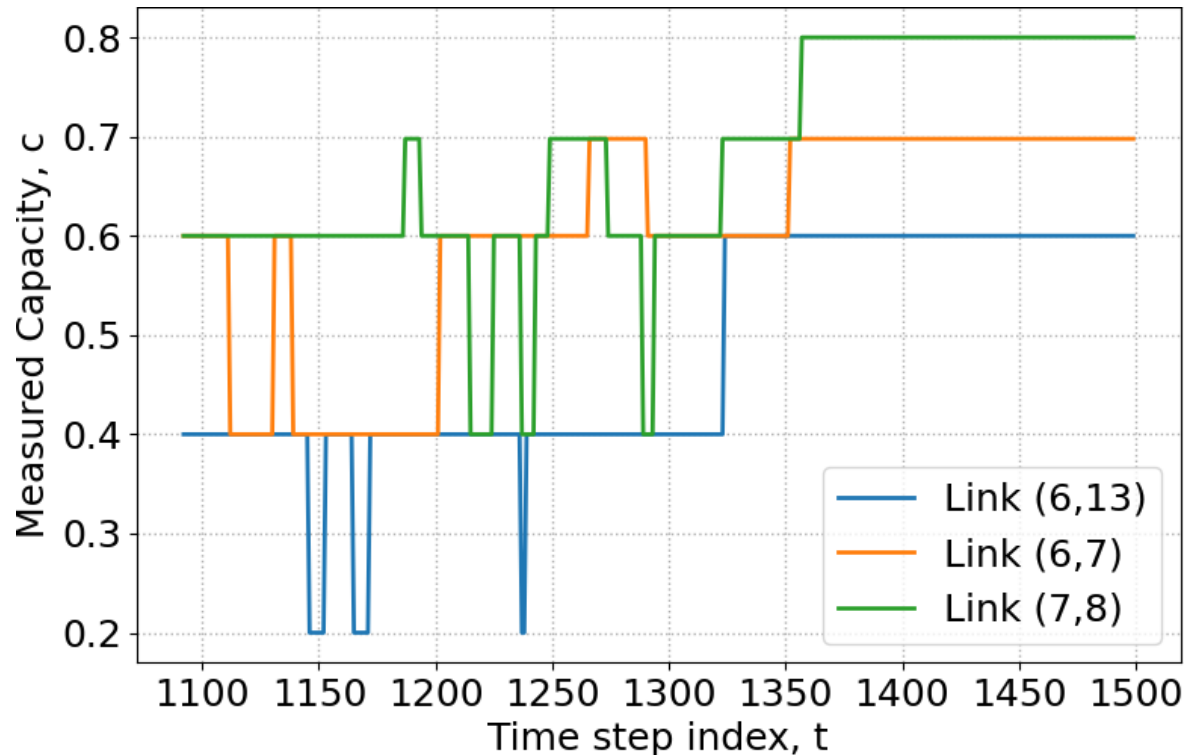
Predictive Network Reconfiguration

- Comparison with two reactive reconfiguration algorithms:
 - Always re-route with Adm Ctrl
 - Never re-route with Adm Ctrl
- Small network with synthetic data and tunable accuracy of attenuation predictions.
- **Large network with measurements is discussed in the next slides...**

Evaluation of PNR using measurements

Challenging scenario with a period of high $c_t^{(k,l)}$ variability due to a rain event.

Node-average admission rates are given by: $\sum_{n=1}^{N-1} z_{n,t} / (N - 1)$



Evaluation of PNR using measurements

Challenging scenario with a period of high $c_t^{(k,l)}$ variability due to a rain event.

Time-average performance gain

$$\frac{\sum_{t=1}^T \sum_{n=1}^{N-1} (z_{n,t}^{MSNR} - z_{n,t}^{React})}{\sum_{t=1}^T \sum_{n=1}^{N-1} z_{n,t}^{React}}$$

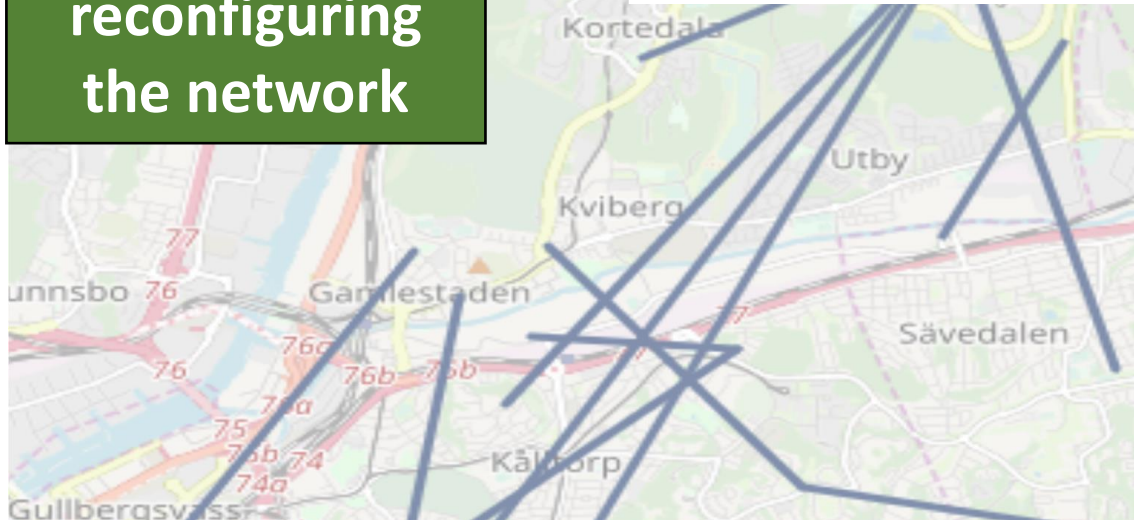
Max. Instant. performance gain

$$\max_{n,t} \left(\frac{z_{n,t}^{MSNR} - z_{n,t}^{React}}{z_{n,t}^{React}} \right)$$

MSNR	Reactive	Time-aver.	Instant.
$H = 2$	ALWAYS	15.49%	263.58%
$H = 2$	NEVER	8.98%	208.01%

Summary

**SDN running
PNR and
reconfiguring
the network**



J. Ostrometzky, G. Zussman, H. Messer, D. Jacoby, and I. Kadota. Predictive Weather-Aware Communication Network Management. US Patent Application No. 17/551,643. December 2021.

(Existing) Data Collection System

**Attenuation Prediction
(AP) Mechanism**

**(Existing) Adaptive Modulation
Mechanism**

**Multi-Step Network Reconfiguration
(MSNR) Algorithm**

**Evaluation of the Predictive Network
Reconfiguration (PNR) framework**